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IBM® Crypto for C version 8.8.1.0

FIPS 140-3 Non-Proprietary Security Policy

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1 General

This document is a non-proprietary FIPS 140-3 Security Policy for the IBM® Crypto for C (ICC) cryptographic module. It contains a specification of the rules under which the module must operate and describes how this module meets the requirements as specified in FIPS PUB 140-3 (Federal Information Processing Standards Publication 140-3) for a security level 1 multi-chip standalone software module.

The table below shows the security level claimed for each of the twelve sections that comprise the FIPS 140-3 standard.

Table 1 - Security Levels

2 Cryptographic Module Specification

The IBM® Crypto for C cryptographic module is implemented in the C programming language. It is packaged as a dynamic (shared) library usable by applications written in a language that supports C language linking conventions (e.g., C, C_{++} , Java, Assembler, etc.) for use on commercially available operating systems. The ICC allows these applications to access cryptographic functions using an Application Programming Interface (API) provided through an ICC import library and based on the API defined by the OpenSSL group.

The software provided to the customer consists of:

- **ICC shared library** (libicclib84.dll for Windows, libicclib084.so for the rest): shared library (executable code) containing proprietary code needed to meet FIPS and functional requirements not provided by OpenSSL (e.g., entropy source, DRBG, self-tests, startup/shutdown), the OpenSSL cryptographic library and the zlib used for entropy estimation. This shared library constitutes the cryptographic module.
- **ICCSIG.txt file**: contains the signature file used for integrity tests.

The cryptographic module takes advantage of the hardware cryptographic acceleration features supported by the testing platforms that are part of the operational environment, as shown in the following table.

The following table presents the operational environments on which version 8.8.1.0 of the cryptographic module was tested and validated. Each operational environment includes the hardware platform, the processor and the operating system. Each row of the table also includes the corresponding version of the module.

#	Operating System	Hardware Platform	Processor	Acceleration
	Red Hat Linux Enterprise Server 8.4 64-bit (Little Endian)	Lenovo ThinkSystem SR630 Intel® Xeon® Gold 5217		AES-NI
	Microsoft Windows Server 2019 64-bit	Lenovo ThinkSystem SR630 Intel® Xeon® Gold 5217		AES-NI
3	Red Hat Linux Enterprise Server 8.4 64-bit (Little Endian) on IBM PowerVM 3.1	IBM Power System S914 (9009-41A)	IBM POWER9	Power ISA
4	Red Hat Linux Enterprise Server 7.9 64-bit (Big Endian) on IBM PowerVM 3.1	IBM Power System S914 (9009-41A)	IBM POWER9	Power ISA
5	IBM AIX 7.2 64-bit (Big Endian) running on IBM PowerVM 3.1	IBM Power System S914 (9009-41A)	IBM POWER9	Power ISA
6	zLinux Red Hat Linux Enterprise Server 8.6 64-bit (Big Endian) on IBM z/VM 7.2	IBM z/15 (8561 T01)	IBM z15	CPACF
7	IBM z/OS 2.3 running on IBM z/VM 7.2	IBM z/15 (8561 T01)	IBM z15	CPACF

Table 2 - Tested Operational Environments

The module maintains its compliance on other operating systems, provided that:

• the operating system meets the operational environment requirements at the module's level of validation;

• the module does not require modification to run in the new environment.

CMVP makes no statement as to the correct operation of the module or the security strengths of the generated keys when so ported if the specific operational environment is not listed on the validation certificate.

The table below lists all approved algorithms of the module, including specific key strengths employed for approved services, and implemented modes of operation. Each algorithm specifies the CAVP certificate for each of the implementations (whether the module was set or unset to take advantage of the processor algorithm acceleration (PAA) or processor algorithm implementation (PAI) capabilities). The selection of the implementations with and without acceleration can be done using the ICC_CAPABILITY_MASK environment variable.

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Table 3 - Approved Algorithms

The Module contains no non-Approved but Allowed security functions, security claimed or otherwise.

The table below lists Non-Approved security functions that are not Allowed in the Approved Mode of Operation.

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Algorithm/Functions	Use/Function		
$SHA-1$	Message digest		
HMAC-MD5	Message authentication code generation, Message authentication code verification		
HMAC-SHA1	Message authentication code generation, Message authentication code verification		
HMAC-DRBG-SHA1	Random number generation		
Hash-DRBG-SHA1	Random number generation		
MDC ₂	Message digest		
RIPEMD	Message digest		
ChaCha20	Symmetric encryption, Symmetric decryption		
ChaCha20-Poly1305	Authenticated encryption, Authenticated decryption		

Table 4 - Non-Approved Not Allowed in the Approved Mode of Operation

The relationship between ICC and IBM applications is shown in the following diagram. ICC comprises a static stub linked into the IBM application which binds the API functions with the shared library containing the cryptographic functionality.

(Figure 1) below depicts the following information:

- **IBM Application** The IBM application using ICC. This contains the application code, and the ICC static stub.
- **IBM Application code** The program using ICC to perform cryptographic functions.
- **ICC static stub**: static library (object code) that is linked into the customer's application and communicates with the Crypto Module. It includes the C headers (source code) containing the API prototypes and other definitions needed for linking the static library. Linked into the calling application to bind the API with the implementation of the cryptographic services in the shared library. This static library is not part of the cryptographic module.
- **ICC shared library** This contains proprietary code needed to meet FIPS requirements and cryptographic services not provided by OpenSSL, a statically linked copy of zlib used by the entropy source for entropy estimation, and a statically linked copy of the OpenSSL cryptographic library.
- **The cryptographic boundary of the cryptographic module** consists of the ICC shared library bounded by the dashed red line in the figure. The signature used for the integrity check of the ICC during its initialization is contained in the file ICCSIG.txt. This file is considered within the cryptographic boundary.
- **The physical perimeter of the operational environment** is defined to be the enclosure of the computer that runs the ICC software.

Figure 1 - Logical Block Diagram

3 Cryptographic Module Ports and Interfaces

The ICC meets the requirements of a multi-chip standalone module. Since the ICC is a software module, its interfaces are defined in terms of the API that it provides. These interfaces are described in the following table¹: Note that because the module is a software only module, there are no physical ports.

Table 5 - Ports and Interfaces

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¹The module does not implement a control output interface.

4 Roles, services, and authentication

The ICC assumes the Crypto-Officer role only (there is no User or Maintenance Role). The module does not support operator identification or authentication. Only a single operator assuming the Crypto Officer role may operate the module at any particular moment in time as concurrent operation is not supported.

The module provides a service indicator that specifies, for a given service, whether the service is approved or non-approved. The module provides the ICC_SetValue() function with the ICC_FIPS_CALLBACK parameter to register a callback function using the following prototype:

void service indicator function(char *function, int nid, int status)

This function is invoked by the module whenever a service is requested, providing the service name (function), the algorithm (nid), and the service indicator (status). A status value of 1 means the service is approved, 0 means non-approved.

The module does not identify nor authenticate any user (in any role) that is accessing the module. The Crypto Officer role is implicitly assumed by the services that are requested. The available services are as follows:

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Table 6 – Roles and Services

The table below lists all approved services that can be used in the approved mode of operation. The abbreviations of the access rights to keys and SSPs have the following interpretation:

- **G** = **Generate**: The module generates or derives the SSP.
- **R** = **Read**: The SSP is read from the module (e.g. the SSP is output).
- **W** = **Write**: The SSP is updated, imported, or written to the module.
- **E** = **Execute**: The module uses the SSP in performing a cryptographic operation.
- **Z** = **Zeroise**: The module zeroises the SSP.

N/A: The service does not access any SSP during its operation.

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Service	Description	Approved Security Functions	Keys and/or SSPs	Role	Access rights to Keys and/or SSPs	Indicator
Show module info	Return module name and versioning information	N/A	None	CO.	N/A	status $=1$
Self-tests	Perform pre-operational and cryptographic algorithm self- tests during power on.	AES. Diffie-Hellman. DSA, EC Diffie- Hellman, ECDSA, DRBG, HKDF, HMAC, RSA, SHS, PBKDF	None	CO.	N/A	status $=1$
On-demand self- tests	Perform cryptographic algorithm self-tests on demand.	AES, Diffie-Hellman, DSA, EC Diffie- Hellman, ECDSA, DRBG, HKDF, HMAC, RSA, SHS, PBKDF	None	CO.	N/A	status $=1$
On-demand integrity test	Perform module integrity test on RSA demand.		None	CO.	N/A	status $=1$
Zeroization	Zeroize SSPs	N/A	All SSPs	CO.	Z	status $=1$
Module linstallation and configuration	Configure module for approved mode of operation	N/A	None	CO	N/A	status $=1$

Table 7 - Approved Services

The following table shows the services and algorithms not allowed in the approved of operation. Requesting these services will implicitly put the module in the non-approved mode of operation.

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Table 8 - Non-Approved Services

5 Software/Firmware security

5.1 Integrity Techniques

The services provided by the Module to a User are effectively delivered using appropriate API calls. When a client process attempts to load an instance of the Module into memory, the Module runs an integrity test and the cryptographic algorithm self-tests. If all the tests pass successfully, the Module makes a transition to the "Operational" state, where the API calls can be used by the client to obtain desired cryptographic services. Otherwise, the Module enters to "Error" state and returns an error to the calling application. When the Module is in "Error" state, no services are available, and all of data input and data output except the status information are inhibited.

The module uses an integrity test which uses a 2048-bit CAVP-validated RSA signature verification (PKCS#1v1.5) and SHA2-256 hashing. This RSA public key is stored inside the shared library.

5.2 On-Demand Integrity Test

Integrity tests are performed as part of the Pre-Operational Self-Tests. They are automatically executed at power-on. Integrity tests can also be requested on demand through the API function ICC_IntegrityCheck.

6 Operational Environment

6.1 Applicability

The IBM® Crypto for C operates in a modifiable operational environment per FIPS 140-3 level 1 specifications. It is part of a commercially available general-purpose operating system executing on the hardware specified in section [2.](#page-4-0)

6.2 Requirements

The following operational rules must be followed by any user of the cryptographic module:

- 1. Since the ICC runs on a general-purpose processor all main data paths of the computer system will contain cryptographic material. The following items need to apply relative to where the ICC will execute:
	- Virtual (paged) memory must be secure (local disk or a secure network)
	- The disk drive where ICC is installed must be in a secure environment.

The above rules must be always upheld in order to ensure continued system security and FIPS 140- 3 approved mode compliance after initial setup of the validated configuration. If the module is removed from the above environment, it is assumed not to be operational in the validated mode until such time as it has been returned to the above environment and re-initialized by the user to the validated condition.

7 Physical Security

The FIPS 140-3 physical security requirements do not apply to the IBM® Crypto for C, since it is a software module.

8 Non-invasive Security

Currently, the non-invasive security is not required by FIPS 140-3 (see NIST SP 800-140F). The requirements of this area are not applicable to the module.

9 Sensitive Security Parameter Management

The following table summarizes the keys and Sensitive Security Parameters (SSPs) that are used by the cryptographic services implemented in the module:

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Table 9 – SSPs

9.1 Random Number Generation

ICC employs a Deterministic Random Bit Generator (DRBG) based on [SP800-90Arev1] for the creation of asymmetric keys. In addition, the module provides a Random Number Generation service to calling applications.

The default algorithm is Hash DRBG using SHA2-256 with no prediction resistance, but another algorithm from the Hash_DRBG, HMAC_DRBG and CTR_DRBG algorithms (see Table [3](#page-14-0) for the complete list) can be also configured.

ICC uses the entropy source to seed the DRBG. The entropy source is a non-physical entropy source ENT (NP) that obtains noise from time jitter produced by the CPU and detected through the CPU high-resolution timer. ENT(NP) is compliant with [SP800-90B], and guarantees an entropy rate of 0.5 bits per bit.

The DRBG entropy input and nonce to form the seed are of the same length (64 bytes $= 512$ bits each) and obtained from separate and independent calls to the entropy source. Then, the DRBG is seeded during initialization with the entropy input and nonce containing 512 bits of entropy ((512 $+$ 512) $*$ 0.5 = 512), and with the entropy input containing 256 bits of entropy (512 $*$ 0.5) during reseeding . Therefore, the DRBG supports 256 bits of effective security strength in its output.

9.2 SSPs Generation

The module generates Keys and SSPs in accordance with FIPS 140-3 IG D.H. The cryptographic module performs Cryptographic Key Generation (CKG) for asymmetric keys as per section 4, example 1, and section 5.1 [SP800-133rev2], compliant with [FIPS186-4] and [SP800-56Arev3]. A seed used for key generation is a direct output from DRBG compliant with [SP800-90A]. The security strength of 256 bits of the DRBG is equal to the security strength of the maximum key size that can be generated by the module.

The key generation services for RSA, Diffie-Hellman and EC key pairs as well as the [SP 800-90A] DRBG have been tested under the CAVP with algorithm certificates found in Table 3.

The ICC provides the following key derivation services in the approved mode of operation:

- PBKDF Key Derivation
	- \circ For PBKDF, the module implements a CAVP compliance tested key derivation function compliant to [SP800-132] and IG D.N. The service returns the key derived from the provided password to the caller.
- HKDF Key Derivation
	- \circ The module provides [SP800-135rev1] compliant key derivation function in accordance with [SP800-56Crev1] two-step key derivation, extraction and expansion procedure. The derived keys provide between 112 and 256 bits of security strength.

9.3 SSPs Establishment

The ICC uses the following key establishment methodologies in the approved mode of operation:

- Diffie-Hellman (DH) shared secret computation
	- \circ The module provides SP800-56Arev3 compliant key agreement schemes according to FIPS 140-3 IG D.F scenario 2 path (1) with DH shared secret computation. The shared secret computation provides between 112 and 200 bits of encryption strength.

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- Elliptic Curve Diffie-Hellman (ECDH) shared secret computation
	- The module provides SP800-56Arev3 compliant key agreement schemes according to FIPS 140-3 IG D.F scenario 2 path (1) with ECDH shared secret computation. The shared secret computation provides between 112 and 256 bits of encryption strength.
- AES key wrapping
	- o The module implements a Key Transport Scheme (KTS) using AES-KW and AES-KWP compliant to [SP800-38F] and IG D.G. The SSP establishment methodology provides between 128 and 256 bits of encryption strength.

9.4 SSPs Import/Export

Keys/SSPs are entered into and output from the ICC module in electronic form through the data input and output interface (i.e. API function parameters). The ICC module does not support manual key entry or intermediate key generation key output. The SSPs are provided to the module via API input parameters in the plaintext form and output via API output parameters in the plaintext form to and from the calling application.

9.5 SSPs Storage

The module does not provide any long-term key storage and no keys are ever stored on the hard disk.

9.6 SSPs Zeroization

The memory occupied by SSPs is allocated by regular memory allocation operating system calls. The calling application that is acting as the Crypto Officer is responsible for calling the appropriate functions provided in the module's API to zeroize the memory areas allocated by the module.

Key zeroization services for cipher contexts are performed via the following API functions.

- ICC BN clear free(): clean up big numbers
- ICC_BN_CTX_free(): clean up memory used by low-level big number arithmetic functions
- ICC_EVP_MD_CTX_cleanup(): clears Message Digest context
- ICC_EVP_CIPHER_CTX_cleanup(): clean up symmetric cipher context
- ICC_RSA_free(): clean up RSA context
- ICC_DSA_free(): clean up DSA context
- ICC DH free(): clean up Diffie-Hellman context
- ICC_EVP_PKEY_free(): clean up asymmetric key contexts
- ICC_HMAC_CTX_free(): clean up HMAC context
- ICC_EC_KEY_free(): clean up ECDSA and ECDH contexts
- ICC_CMAC_CTX_free(): clean up CMAC context
- ICC_AES_GCM_CTX_free(): clean up AES-GCM context
- ICC_RNG_CTX_free(): clean up RNG context

The zeroization functions overwrite the memory occupied by SSPs with "zeros" and deallocate the memory with the regular memory deallocation operating system call. The completion of a zeroization routine(s) will indicate that a zeroization procedure succeeded.

10 Self-tests

The ICC module implements a number of self-tests to check proper functioning of the module. This includes pre-operational self-tests and conditional self-tests.

Pre-operational integrity test and Cryptographic Algorithm Self-Tests (CASTs) are automatically invoked by the module when the module is powered on from the default entry point (DEP) of the shared library..

When the module is performing self-tests, no API functions are available, and no data output is possible until the self-tests are successfully completed. After the pre-operational self-tests and CASTs are successfully completed, the module turns to approved mode of operation. Requesting any services from Table 8 will implicitly put the module in the non-approved mode of operation.

The module performs self-tests automatically when it is loaded. Self-tests can also be requested on demand through the API functions ICC_SelfTest() and ICC_IntegrityCheck().

Whenever the startup tests are initiated the module performs the following; if any of these tests fail, the module enters the error state:

10.1 Pre-operational Software Integrity Test

The module performs a pre-operational software integrity test automatically when the module is powered on, before the module transitions into the operational state. The integrity test is performed with a 2048-bit CAVP-validated RSA signature verification (PKCS#1v1.5) and SHA2-256 hashing. This RSA public key is stored inside the shared library.

Prior to the invocation of the integrity test, the module runs the conditional Cryptographic Algorithm Self-Test (CAST) for RSA (2048-bit keys with SHA2-256) which verifies the proper functioning of all algorithms used as part of the integrity test.

10.2 Conditional Self-Tests

The following sections describe the conditional tests supported by the IBM® Crypto for C.

10.2.1 Cryptographic Algorithm Self-Tests

The IBM® Crypto for C runs all Cryptographic Algorithm Self-Tests during power-up, and consequently before the first operational use of the cryptographic algorithms. These tests are detailed in the following table.

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Cryptographic Algorithm	Notes		
RSA with 2048-bit keys and SHA2-256	Separate signature generation/ verification KAT are performed		
ECDSA with curves P-384 and B-233 and using SHA2-256	Separate signature generation / verification KAT are performed		
DSA with L=2048, N=224 and SHA2-256	Signature verification KAT		
Hash DRBG with SHA-224, SHA-256, SHA-384 and SHA-512 HMAC DRBG with SHA-224, SHA-256, SHA-384 and SHA-512 CTR DRBG with AES-128, AES-192 and AES-256	Each DRBG mode tested separately.		
DRBG health tests	Health tests according to section 11.3 of [SP800- 90Arev11		
HKDF using SHA2-256	KAT		
PBKDF using SHA2-256	KAT		
Diffie-Hellman "Z" computation with 2048-bit key KAT			
EC Diffie-Hellman "Z" computation with P-521 curve	KAT		
Repetitive Counter Test (RCT)	Startup tests of the ENT(NP) entropy source. Performed on 1024 consecutive samples.		
Adaptive Proportion Test (APT)	Startup tests of the ENT(NP) entropy source. Performed on 1024 consecutive samples.		

Table 11 - Cryptographic Algorithm Self-Tests

10.2.2 Pairwise Consistency Test

The IBM® Crypto for C does generate asymmetric keys and performs all required pair-wise consistency tests. The consistency of the keys is tested by the calculation and verification of a digital signature. If the digital signature cannot be verified, the test fails. Pair-wise consistency tests are performed on the following algorithms:

- ECDSA signature generation and verification using SHA2-256.
- RSA signature generation and verification using SHA2-256.
- Diffie-Hellman according to section 5.6.2.1.4 of [SP800-56Arev3].

EC Diffie-Hellman is covered by ECDSA PCT as allowed by IG 10.3, additional comment 1.

10.2.3 Entropy Health Test

The ICC module performs health tests during the startup of the ENT(NP), and continuously during its operation, to detect intermittent and permanent failures in the noise source. The health tests implemented are the Repetitive Count Test (RCT) and Adaptive Proportion Test (APT), both compliant with the requirements of SP800-90B, and the minimum-entropy assessment test, which analyzes whether the noise source provides the expected entropy rate using the min-entropy calculation formula as specified in section 2.1 of SP800-90B.

If the ICC module detects a permanent failure in any of the health tests, the module transitions to the error state and an error message is shown ("Insufficient entropy").

10.3 Error Handling

When errors are detected (e.g., self-test failure) then all security related functions are disabled and no partial data is exposed through the data output interface. The only way to transition from the error state to an operational state is to reinitialize the cryptographic module (from an uninitialized state). The error state can be retrieved via the Show Status service.

11 Life-cycle assurance

11.1 Delivery and Operation

The following steps must be performed to install and initialize the module for operating in a FIPS 140-3 compliant manner:

- 1. The operating system must be configured to operate securely and to prevent remote login. This is accomplished by disabling all services (within the Administrative tools) that provide remote access (e.g., – ftp, telnet, ssh, and server) and disallowing multiple operators to log in at once.
- 2. Before the module initialization, the user has a choice to configure the default DRBG algorithm to use. This can be set using the environment variable 'ICC_RANDOM_GENERATOR'.
- *3.* The module is initialized automatically when the shared library is loaded in the calling application process space. The module executes the pre-operational self tests (POST) and, if they are successful, the module enters the approved mode of operation. The calling application must include the following calling sequence to have access to the cryptographic services:
	- **ICC Init()** creates the crypto module context.
	- **ICC Attach()** binds the cryptographic functions with the API entry points.

11.2 Crypto Officer Guidance

It is the responsibility of the Crypto-Officer to configure the operating system to operate securely.

The services provided by the Module to a User are effectively delivered by using the appropriate API calls. When a client process attempts to load an instance of the Module into memory, the Module runs an integrity test and several of cryptographic functionality self-tests. If all the tests pass successfully, the Module makes a transition to the "Operational" state, where the API calls can be used by the client to obtain desired cryptographic services. Otherwise, the Module enters to "Error" state and returns an error to the calling application. When the Module is in "Error" state, no services are available, and all of data input and data output except the status information are inhibited.

The Crypto Officer shall consider the following requirements and restrictions when using the module:

- 1. The AES algorithm in XTS mode can be only used for the cryptographic protection of data on storage devices, as specified in [SP800-38E]. The length of a single data unit encrypted with the XTS-AES shall not exceed 2²⁰ AES blocks (16MB of data).
- 2. To meet the requirement in [FIPS140-3-IG] C.I, the module implements a check to ensure that the two AES keys used in the XTS-AES algorithm are not identical.
- 3. AES-GCM IV is constructed in compliance with IG C.H scenario 1. In case the module's power is lost and then restored, the keys used for the AES GCM encryption/decryption shall be re-distributed. The GCM is used in the context of TLS version 1.2. The mechanism for IV generation is compliant with RFC 5288 as described in Section 3.3.1 of SP800-52rev2. The design of the TLS protocol implicitly ensures that the nonce explicit, or counter portion of the IV will not exhaust all its possible values.
- 4. The module also offers an AES-GCM implementation under the context of Scenario 5 of IG C.H. The protocol that provides this compliance is TLS 1.3, using the ciphersuites that

explicitly select AES-GCM as the encryption/decryption cipher. The module supports acceptable AES-GCM ciphersuites from Section 3.3.1 of SP800-52rev2.

The design of the TLS protocol implicitly ensures that the nonce explicit, or counter portion of the IV will not exhaust all its possible values. In the event the module's power is lost and restored, the consuming application must ensure that new AES-GCM keys encryption or decryption under this scenario are established. TLS 1.3 provides session resumption, but the resumption procedure derives new AES-GCM encryption keys.

5. For PBKDF, the module implements a CAVP compliance tested key derivation function compliant to [SP800-132] and IG D.N. The service returns the key derived from the provided password to the caller.

PBKDF is implemented to support the option 1a specified in section 5.4 of [SP800-132]. The keys derived from [SP800-132] map to section 4.1 of [SP800-133rev2] as indirect generation from DRBG.

In accordance with [SP800-132], the following requirements shall be met:

- a. Derived keys shall only be used in storage applications. The Master Key (MK) shall not be used for other purposes. The length of the MK or Data Protection Key (DPK) shall be of 112 bits or more.
- b. A portion of the salt, with a length of at least 128 bits, shall be generated randomly using the SP800-90A DRBG,
- c. The iteration count shall be equal or greater than 1000, so as to make the key derivation computationally intensive.
- d. Passwords or passphrases, used as an input for the PBKDF, shall not be used as cryptographic keys.
- e. The length of the password or passphrase shall be at least ten characters long, and may consist of lower-case, upper-case, numeric, or special characters. At a minimum length of ten characters, and assuming a worst case scenario where the password uses a combination of only lower case and numbers (36 symbols), the chance of randomly quessing this password is $1/36^{10} = 3.656 10^{-15}$.
- 6. For SHA-3 algorithms, the module implements HMAC with SHA3-224, SHA3-256, SHA3-384, SHA3-512. The CAVP certificates have been obtained for the HMAC and HKDF algorithms as well as for all the SHA-3 implementations. The CAVP certificates are listed in Table 3 in Section 2.
- 7. The module implements FIPS 186-4 RSA SigGen and SigVer. RSA SigGen is supported with key sizes of 2048, 3072, 4096 bits while RSA SigVer is supported with 1024, 2048, 3072, 4096 bits. All RSA key sizes have been CAVP tested with the certificates listed in Table 3 in Section 2.
- 8. For Diffie-Hellman or EC Diffie-Hellman shared secret computation, the module has to comply with the assurances found in Section 5.6.2 of [SP800-56Arev3] and IG D.F. The operator must obtain the ephemeral Diffie-Hellman or EC Diffie-Hellman key pairs on both ends either by using the approved key pair generation service provided by the module, or by using another FIPS-validated module. As part of the key pair generation service, the module internally performs the full key validation of the generated key pair. Similarly, the shared secret computation service internally performs the full public key validation of the peer public key, complying with Sections 5.6.2.2.1 and 5.6.2.2.2 of [SP800-56Arev3].

The module code is a component provided to IBM products, not a product on its own. Typically it is provided as part of IBM's SSL component and creates packaging with the OS specific install tools.

The module's End-of-Life/sanitization procedure can take one of two forms:

- OS uninstall which removes the package after checking that no currently installed package still depends on it. The module does not possess persistent storage of SSPs. The SSP value only exists in volatile memory and that value vanishes when the module is powered off. The procedure for secure sanitization of the module at the end of life is simply to power it off, which is the action of zeroization of the SSPs . As a result of this sanitization via power-off, the SSP is removed from the module, so that the module may either be distributed to other operators or disposed.
- Alternatively the package may be upgraded and replaced by a newer version.

12 Mitigation of other attacks

The cryptographic module is not designed to mitigate any specific attacks.

Appendix A. Glossary and Abbreviations

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- **RNG** Random Number Generator
- **RSA** Rivest, Shamir, Addleman
- **SHA** Secure Hash Algorithm
- **SHS** Secure Hash Standard
- **SSH** Secure Shell
- **TDES** Triple-DES
- **XTS** XEX-based Tweaked-codebook mode with cipher text Stealing

Appendix B. References

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<https://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP.800-133r2.pdf>